

## Algebraic Coding Theory e-Summer School **ACT21**

# An Orbital Construction of Optimum Distance Flag Codes

Miguel Ángel Navarro-Pérez

Joint work with Clementa Alonso-González and Xaro Soler-Escrivà

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# Outline

- 1 Constant dimension codes
- 2 Flag codes
  - Optimum distance full flag codes
- 3 Our construction

# Notation

Let...

- $q$  be a prime power,
- $\mathbb{F}_q$  denote the *finite field* with  $q$  elements,
- $n$  be a positive integer.
- $\mathbb{F}_q^n$  is the  $n$ -dimensional vector space over  $\mathbb{F}_q$ ,
- $\mathcal{G}_q(k, n)$ , the *Grassmannian* of dimension  $k$ .

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# Constant dimension codes

## Definition

A **constant dimension code**  $\mathcal{C}$  in  $\mathcal{G}_q(k, n)$  is a nonempty collection of  $k$ -dimensional subspaces of  $\mathbb{F}_q^n$ .

Given two subspaces  $\mathcal{U}, \mathcal{V}$  of  $\mathbb{F}_q^n$ , with  $\dim(\mathcal{U}) = \dim(\mathcal{V}) = k$ , then

$$d_S(\mathcal{U}, \mathcal{V}) = 2(k - \dim(\mathcal{U} \cap \mathcal{V})).$$

## Group action

For every  $1 \leq k \leq n - 1$ , the group  $GL(n, q)$  acts on the Grassmannian  $\mathcal{G}_q(k, n)$ :

$$\mathcal{G}_q(k, n) \times GL(n, q) \rightarrow \mathcal{G}_q(k, n), \quad \mathcal{V} \cdot A \mapsto \text{rowsp}(VA),$$

with  $V$  a generator matrix of  $\mathcal{V}$ .

## Definition

Given a subspace  $\mathcal{V} \in \mathcal{G}_q(k, n)$  and a subgroup  $T$  of  $GL(n, q)$ , the **orbit code** generated by  $\mathcal{V}$  (under the action of  $T$ ) is

$$\text{Orb}_T(\mathcal{V}) = \{\mathcal{V} \cdot A \mid A \in T\}.$$

Its stabilizer is the subgroup  $\text{Stab}_T(\mathcal{V}) = \{A \in T \mid \mathcal{V} \cdot A = \mathcal{V}\}$  and it holds  $|\text{Orb}_T(\mathcal{V})| = |T|/|\text{Stab}_T(\mathcal{V})|$ .

# Spread codes

## Definition

A  **$k$ -spread code**  $\mathcal{S}$  of  $\mathbb{F}_q^n$  is a constant dimension code in  $\mathcal{G}_q(k, n)$  such that every nonzero vector of  $\mathbb{F}_q^n$  lies on one, and only one, subspace in  $\mathcal{S}$ . **Vector space partition**

## Properties:

- $k$ -spread codes exist  $\Leftrightarrow k$  divides  $n$ ,
- $d_{\mathcal{S}}(\mathcal{S}) = 2k$  and
- $|\mathcal{S}| = \frac{q^n - 1}{q^k - 1}$ .

## Spreads codes

In [13], Manganiello, Gorla and Rosenthal present the next construction of a  $k$ -spread code.

Theorem ([13, Th. 1])

Let  $P$  be the companion matrix of a primitive polynomial of degree  $k$  in  $\mathbb{F}_q[x]$ . Then the set

$$\mathcal{S} = \{\text{rowsp}(I_k | P^i), \text{rowsp}(0_k | I_k) \mid 1 \leq i \leq q^k - 2\} \quad (1)$$

is a  $k$ -spread of  $\mathbb{F}_q^{2k}$ .

In [16], the authors provide an orbital description of  $\mathcal{S}$ .



## Spreads as orbit codes

Consider matrices in  $GL(2k, q)$

$$g = \left[ \begin{array}{c|c} 0_k & I_k \\ \hline I_k & 0_k \end{array} \right] \quad \text{and} \quad g_i = \left[ \begin{array}{c|c} I_k & P^i \\ \hline 0_k & I_k \end{array} \right], \quad 0 \leq i \leq q^k - 2.$$

Form the group

$$G = \langle g, g_i, \mid 0 \leq i \leq q^k - 2 \rangle.$$

and take  $\mathcal{U}_k = \text{rowsp}(I_k \mid 0_k)$ .

Lemma ([16, Lemma 15])

The orbit code  $\text{Orb}_G(\mathcal{U}_k) = \mathcal{S}$  is the  $k$ -spread code in (1).

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# Full flags

## Definition

A **full flag** on  $\mathbb{F}_q^n$  is a sequence

$$\mathcal{F} = (\mathcal{F}_1, \dots, \mathcal{F}_{n-1})$$

such that

- $\mathcal{F}_1 \subsetneq \mathcal{F}_2 \subsetneq \dots \subsetneq \mathcal{F}_{n-1}$ .
- $\dim(\mathcal{F}_i) = i$ ,  **$i$ -th subspace of  $\mathcal{F}$ .**

## Full flag codes

The **flag distance** between two full flags  $\mathcal{F}$  and  $\mathcal{F}'$  on  $\mathbb{F}_q^n$  is

$$d_f(\mathcal{F}, \mathcal{F}') = \sum_{i=1}^{n-1} d_S(\mathcal{F}_i, \mathcal{F}'_i).$$

## Definition

A **full flag code** on  $\mathbb{F}_q^n$  is a nonempty collection of full flags. Its **minimum distance** is

$$d_f(\mathcal{C}) = \min\{d_f(\mathcal{F}, \mathcal{F}') \mid \mathcal{F}, \mathcal{F}' \in \mathcal{C}, \mathcal{F} \neq \mathcal{F}'\}$$

# Optimum distance full flag codes

## Definition

A full flag code  $\mathcal{C}$  on  $\mathbb{F}_q^n$  is said to be an **optimum distance flag code** if

$$d_f(\mathcal{C}) = \begin{cases} \frac{n^2}{2} & \text{if } n \text{ is even,} \\ \frac{n^2-1}{2} & \text{if } n \text{ is odd.} \end{cases}$$

# Projected codes of flag code

## Definition

Let  $\mathcal{C}$  be a full flag code on  $\mathbb{F}_q^n$ . For every  $i \in \{1, \dots, n-1\}$ , the  **$i$ -projected code of  $\mathcal{C}$**  is the subspace code  $\mathcal{C}_i$  given by the **set of all the  $i$ -th subspaces** of flags in  $\mathcal{C}$ .

$$\mathcal{C}_i = \{\mathcal{F}_i \in \mathcal{G}_q(i, n) \mid (\mathcal{F}_1, \dots, \mathcal{F}_i, \dots, \mathcal{F}_{n-1}) \in \mathcal{C}\}.$$

- $\mathcal{C}_i \subseteq \mathcal{G}_q(i, n)$ ,
- $|\mathcal{C}_i| \leq |\mathcal{C}|$ .
- If  $|\mathcal{C}_1| = \dots = |\mathcal{C}_{n-1}| = |\mathcal{C}|$  we say that  $\mathcal{C}$  is **disjoint**.

# Optimum distance flag codes

## Theorem ([3, Theorem 3.11])

Let  $\mathcal{C}$  be a full flag code on  $\mathbb{F}_q^n$ . The following statements are equivalent.

- 1  $\mathcal{C}$  is an **optimum distance** flag code.
- 2  $\mathcal{C}$  is disjoint and every  $\mathcal{C}_i$  attains the maximum (subspace) distance.

For  $n = 2k$ , we have the following result.

## Theorem ([3, Th. 3.12])

Let  $\mathcal{C}$  be an optimum distance full flag code on  $\mathbb{F}_q^{2k}$ . Then  $|\mathcal{C}| \leq q^k + 1$  and the equality holds if, and only if,  $\mathcal{C}_k$  is a  $k$ -spread of  $\mathbb{F}_q^{2k}$ .

# Orbit flag codes

Consider a full flag  $\mathcal{F}$  on  $\mathbb{F}_q^n$  and a matrix  $A \in GL(n, q)$ , then

$$\mathcal{F} \cdot A = (\mathcal{F}_1 \cdot A, \dots, \mathcal{F}_{n-1} \cdot A)$$

is a full flag.

## Definition

Given a group  $T \subseteq GL(n, q)$ , the **orbit flag code** generated by  $\mathcal{F}$  (under the action of  $T$ ) is  $\text{Orb}_T(\mathcal{F})$ .

- $\text{Orb}_T(\mathcal{F})_i = \text{Orb}_T(\mathcal{F}_i)$
- $\text{Stab}_T(\mathcal{F}) = \bigcap_{i=1}^{n-1} \text{Stab}_T(\mathcal{F}_i)$ .



# Orbit flag codes

Proposition ([2, Prop. 3.5])

The code  $\text{Orb}_T(\mathcal{F})$  is disjoint iff  $\text{Stab}_T(\mathcal{F}_1) = \cdots = \text{Stab}_T(\mathcal{F}_{n-1})$ .

Proposition ([2, Prop. 3.6])

Take  $n = 2k$ . If  $d_S(\text{Orb}_T(\mathcal{F}_k)) = 2k$ , then  $\text{Stab}_T(\mathcal{F}_i) \subseteq \text{Stab}_T(\mathcal{F}_k)$ , for all  $1 \leq i \leq n - 1$ .

Proposition ([2, Prop. 3.7])

Take  $n = 2k$ . If  $d_S(\text{Orb}_T(\mathcal{F}_k)) = 2k$  and  $\text{Stab}_T(\mathcal{F}_k) \subseteq \text{Stab}_T(\mathcal{F}_i)$ , then  $\text{Orb}_T(\mathcal{F})$  is an optimum distance flag code.

In particular, if  $\text{Stab}_T(\mathcal{F}_k) = \{I_n\}$ , we are done.

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Starting from a  $k$ -spread

Let  $P$  be the companion matrix of a primitive polynomial of degree  $k$  in  $\mathbb{F}_q[x]$  and consider the  $k$ -spread

$$\mathcal{S} = \{\text{rowsp}(I_k | P^i), \text{rowsp}(0_k | I_k) \mid 1 \leq i \leq q^k - 2\}. \quad (2)$$

Recall that  $\mathcal{S} = \text{Orb}_G(\text{rowsp}(I_k | 0_k))$  with

$$G = \left\langle \left[ \begin{array}{c|c} 0_k & I_k \\ \hline I_k & 0_k \end{array} \right], \left[ \begin{array}{c|c} I_k & P^i \\ \hline 0_k & I_k \end{array} \right] \mid 0 \leq i \leq q^k - 2 \right\rangle.$$

- We start analyzing the group structure of  $G$ . Consider a primitive element  $\alpha \in \mathbb{F}_{q^k}$ .
- The field isomorphism  $\phi : \mathbb{F}_{q^k} \rightarrow \mathbb{F}_q[P]$  such that  $\alpha \mapsto \phi(\alpha) = P$  induces a group isomorphism between

$$\bar{G} = \left\langle \left[ \begin{array}{cc} 0 & 1 \\ 1 & 0 \end{array} \right], \left[ \begin{array}{cc} 1 & \alpha^i \\ 0 & 1 \end{array} \right], 1 \leq i \leq q^k - 2 \right\rangle$$

and

$$G = \left\langle \left[ \begin{array}{c|c} 0_k & I_k \\ \hline I_k & 0_k \end{array} \right], \left[ \begin{array}{c|c} I_k & P^i \\ \hline 0_k & I_k \end{array} \right], 1 \leq i \leq q^k - 2 \right\rangle.$$

Proposition ([2, Prop. 4.2])

$$G \cong \bar{G} = \begin{cases} SL(2, q^k) & \text{if } \text{char}(\mathbb{F}_q) = 2, \\ SL(2, q^k) \rtimes \left\langle \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \right\rangle, & \text{otherwise.} \end{cases}$$

Remark:

Moreover, we can restrict ourselves to the subgroup  $\psi(SL(2, q^k))$  of  $G$  since

$$\mathcal{S} = \text{Orb}_{\psi(SL(2, q^k))}(\text{rowsp}(I_k | 0_k)).$$

Recall that, if we find a subgroup  $T$  of  $\psi(SL(2, q^k))$  such that:

- $\mathcal{S} = \text{Orb}_T(\mathcal{V})$  and
- $\text{Stab}_T(\mathcal{V}) = \{I_n\}$ ,

then we can easily provide an optimum distance flag code. In that case,  $|T| = q^k + 1 = |\mathcal{S}|$ .

Let  $\bar{H}$  be a Singer subgroup of  $SL(2, q^k)$  and consider  $H = \psi(\bar{H})$ . Then  $H$  is cyclic of order  $q^k + 1$  and, for every  $\mathcal{V} \in \mathcal{S}$ , it holds:

Lemma ([2, Lemma 4.11])

$$\text{Stab}_H(\mathcal{V}) = \begin{cases} \{I_n\} & \text{if } \text{char}(\mathbb{F}_q) = 2, \\ \{I_n, -I_n\} & \text{if } \text{char}(\mathbb{F}_q) \neq 2. \end{cases}$$

Theorem ([2, Th. 4.14])

Let  $\mathcal{F}$  be a full flag on  $\mathbb{F}_q^{2k}$  with  $\mathcal{F}_k \in \mathcal{S}$ . Then  $\text{Orb}_H(\mathcal{F})$  is an optimum distance flag code with cardinality

$$\begin{cases} q^k + 1 & \text{if } \text{char}(\mathbb{F}_q) = 2, \\ \frac{q^k + 1}{2} & \text{if } \text{char}(\mathbb{F}_q) \neq 2. \end{cases}$$

## Proposition ([2, Prop. 4.15])

For odd characteristic: Let  $\mathcal{F}$  and  $\mathcal{F}'$  be full flags on  $\mathbb{F}_q^{2k}$  with  $\mathcal{F}_k, \mathcal{F}'_k \in \mathcal{S}$  such that  $\text{Orb}_H(\mathcal{F}_k) \neq \text{Orb}_H(\mathcal{F}'_k)$ . Then

$$\mathcal{C} = \text{Orb}_H(\mathcal{F}) \cup \text{Orb}_H(\mathcal{F}')$$

is an optimum distance flag code with size  $q^k + 1$ .



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