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Construction of an optimal convolutional code of rate $1/2$

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July 11, 2022

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Introduction

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- Maximum Distance Profile (MDP) codes
- **PROBLEM:** Build convolutional codes with optimal distances

Main goal

Our main goal with this research was:

- 1 propose a new construction of convolutional codes. In particular, build convolutional codes with maximum free distance

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Two generator matrices $G_1(D), G_2(D) \in \mathbb{F}_q[D]^{k \times n}$ are said to be **equivalent generator matrices** if

$$\text{Im}_{\mathbb{F}_q[D]} G_1(D) = \text{Im}_{\mathbb{F}_q[D]} G_2(D),$$

which happens if and only if $G_1(D) = U(D)G_2(D)$ for some unimodular matrix $U(D) \in \mathbb{F}_q[D]^{k \times k}$.

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- The **degree** of a code is the external degree of a minimal generator matrix.

The **free distance** of a convolutional code is defined as:

$$d_{free}(\mathcal{C}) = \min\{wt(v(D)) \mid v(D) \in \mathcal{C}, v(D) \neq 0\},$$

where $wt(v(D))$ is the Hamming weight of a polynomial vector

$v(D) = \sum_{t=0}^{\deg(v(D))} v_t D^t \in \mathbb{F}_q[D]^n$ that is defined as

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Maximum distance separable (MDS) codes have maximum free distance in the class of convolutional codes of rate k/n and degree δ , i.e., are the ones with free distance equal to the generalized **Singleton bound** (upper bound)

$$(n-k) \left(\left\lfloor \frac{\delta}{k} \right\rfloor + 1 \right) + \delta + 1.$$

Main idea

Main idea

Consider a generator matrix $G(D)$ of a convolutional code over a field \mathbb{F}_q and construct a generator matrix of a new convolutional code over the same field \mathbb{F}_q and of higher degree, by repeating the coefficients of $G(D)$ of certain degrees.

Advantage

- Build a convolutional code MDS defined over a smaller field than the existing constructions of MDS convolutional codes with the same degree.

Justesen's Theorem

Theorem ([4])

For $n \geq 2$ and $|\mathbb{F}_q| \geq n + 1$, set $s_j := \lceil (j - 1)(|\mathbb{F}_q| - 1)/n \rceil$ for $j = 2, \dots, n$ and

$$\delta := \begin{cases} \lfloor \frac{2}{9} |\mathbb{F}_q| \rfloor & \text{if } n = 2, \\ \lfloor \frac{1}{3} |\mathbb{F}_q| \rfloor & \text{if } 3 \leq n \leq 5, \\ \lfloor \frac{1}{2} |\mathbb{F}_q| \rfloor & \text{if } n \geq 6 \end{cases}$$

Moreover, let α be a primitive element of \mathbb{F}_q , and set

$$g_1(x) := \prod_{k=1}^{\delta} (x - \alpha^k),$$

$$g_j(x) := g_1(x\alpha^{-s_j})$$

for $j = 2, \dots, n$. Then

$$G(D) = [g_1(D) \cdots g_n(D)]$$

is the generator matrix of an MDS $(n, 1, \delta)$ convolutional code with free distance equal to $n(\delta + 1)$.

New construction

Let $n = 2$ and $\delta = 2$

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Construction Idea

Let α be a primitive element of \mathbb{F}_q , with $q = 11, 13$. Then

$$\tilde{G}(D) = G_0 + G_1 D + G_2 D^2 + G_2 D^3 + G_1 D^4 + G_0 D^5, \text{ with}$$

$$G_0 = \begin{bmatrix} \alpha^3 & \alpha^3 \end{bmatrix}$$

$$G_1 = \begin{bmatrix} -\alpha^2 - \alpha & \alpha^{-s_j}(-\alpha^2 - \alpha) \end{bmatrix}$$

$$G_2 = \begin{bmatrix} 1 & 1 \end{bmatrix}$$

is the generator matrix of a $(2, 1, 5)$ MDS convolutional code.

Primitive elements (α):

- \mathbb{F}_{11} : 2, 6, 7, 8
- \mathbb{F}_{13} : 2, 6, 7, 11

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- \mathbb{F}_{13} : 2, 6, 7, 11

Theorem (Construction)

Let $\alpha = 2$ be a primitive element of \mathbb{F}_{11} . Then

$$\tilde{G}(D) = G_0 + G_1 D + G_2 D^2 + G_2 D^3 + G_1 D^4 + G_0 D^5,$$

with $G_0 = \begin{bmatrix} 8 & 8 \end{bmatrix}$, $G_1 = \begin{bmatrix} 5 & 6 \end{bmatrix}$ and $G_2 = \begin{bmatrix} 1 & 1 \end{bmatrix}$ is the generator matrix of a $(2, 1, 5)$ MDS convolutional code.

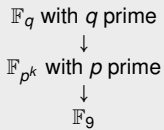
Theorem (Construction)

Let $\alpha = 6$ be a primitive element of \mathbb{F}_{11} . Then

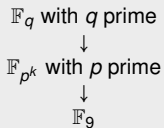
$$\tilde{G}(D) = G_0 + G_1 D + G_2 D^2 + G_2 D^3 + G_1 D^4 + G_0 D^5,$$

with $G_0 = \begin{bmatrix} 7 & 7 \end{bmatrix}$, $G_1 = \begin{bmatrix} 2 & 9 \end{bmatrix}$ and $G_2 = \begin{bmatrix} 1 & 1 \end{bmatrix}$ is the generator matrix of a $(2, 1, 5)$ MDS convolutional code.

Construction Idea



Construction Idea



$$\mathbb{F}_9 = \{0, 1, 2, \alpha, \alpha + 1, \alpha + 2, 2\alpha, 2\alpha + 1, 2\alpha + 2\}$$








- Primitive elements: $\alpha + 1, \alpha + 2, 2\alpha + 1, 2\alpha + 2$

Future work

Parameters

- δ
- n

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Thank you for your attention

Acknowledgments

This work was supported by The Center for Research and Development in Mathematics and Applications (CIDMA) through the Portuguese Foundation for Science and Technology (FCT - Fundação para a Ciência e a Tecnologia), UIDB/04106/2020 and UIDP/04106/2020. Zita Abreu is also supported by FCT grant UI/BD/151186/2021.



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