

The Density of Extremal Codes With Sublinearity

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July 12th, 2022

Joachim Rosenthal's 60th birthday

Joint work with Anina Gruica, Anna-Lena Horlemann and Alberto Ravagnani



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$$\delta_q^D(\mathbb{F}_{q^m}^n, S, 0, d) := \frac{|\{\mathcal{C} \subseteq \mathbb{F}_{q^m}^n : |\mathcal{C}| = S, D(\mathcal{C}) \geq d\}|}{|\{\mathcal{C} \subseteq \mathbb{F}_{q^m}^n : |\mathcal{C}| = S\}|}$$

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$$\delta_q^D(\mathbb{F}_{q^m}^n, S, \ell, d) := \frac{|\{\mathcal{C} \subseteq \mathbb{F}_{q^m}^n : |\mathcal{C}| = S, D(\mathcal{C}) \geq d, \mathcal{C} \text{ is } \mathbb{F}_{q^\ell}\text{-linear}\}|}{|\{\mathcal{C} \subseteq \mathbb{F}_{q^m}^n : |\mathcal{C}| = S, \mathcal{C} \text{ is } \mathbb{F}_{q^\ell}\text{-linear}\}|}$$

the density of \mathbb{F}_{q^ℓ} -linear codes with $D(\mathcal{C}) \geq d$.

Theorem

Let $(S_q)_{q \in Q}$ be a sequence of integers with $S_q \geq 4$ for all $q \in Q$ and for which $\lim_{q \rightarrow +\infty} S_q$ exists. We fix $1 \leq d \leq \max\{D(x, y) : x, y \in \mathbb{F}_q^n\}$ and let $(v_q^D(\mathbb{F}_q^n, d-1))_{q \in Q}$ be a sequence where each $v_q^D(\mathbb{F}_q^n, d-1)$ describes the size of the ball of radius $d-1$ in \mathbb{F}_q^n . We have

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(i)

$$\max \left\{ \liminf_{q \rightarrow +\infty} \left(1 - \frac{\mathbf{v}_q^D(\mathbb{F}_q^n, d-1) S_q^2}{2q^n} \right), 0 \right\} \leq \liminf_{q \rightarrow +\infty} \delta_q^D(\mathbb{F}_q^n, S_q, 0, d).$$

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In particular, if $S_q \in o\left(\sqrt{\frac{q^n}{\mathbf{v}_q^D(\mathbb{F}_q^n, d-1)}}\right)$ as $q \rightarrow +\infty$, then

$$\lim_{q \rightarrow +\infty} \delta_q^D(\mathbb{F}_q^n, S_q, 0, d) = 1.$$

i.e. the class of non-linear codes with $D(\mathcal{C}) \geq d$ is **dense**.

Theorem (cont.)

(ii) If $S_q \in \Omega \left(\sqrt{\frac{q^n}{\mathbf{v}_q^D(\mathbb{F}_q^n, d-1)}} \right)$ as $q \rightarrow +\infty$, then

$$\limsup_{q \rightarrow +\infty} \delta_q^D(\mathbb{F}_q^n, S_q, 0, d) \leq \limsup_{q \rightarrow +\infty} \left(1 - \frac{\mathbf{v}_q^D(\mathbb{F}_q^n, d-1) S_q^2}{2q^n + \mathbf{v}_q^D(\mathbb{F}_q^n, d-1) S_q^2} \right).$$

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In particular, if $S_q \in \omega \left(\sqrt{\frac{q^n}{\mathbf{v}_q^D(\mathbb{F}_q^n, d-1)}} \right)$ as $q \rightarrow +\infty$, then

$$\lim_{q \rightarrow +\infty} \delta_q^D(\mathbb{F}_q^n, S_q, 0, d) = 0,$$

i.e. the class of non-linear codes with $D(\mathcal{C}) \geq d$ is **sparse**.

Theorem

We fix $1 \leq d \leq \max\{D(x, y) : x, y \in \mathbb{F}_{q^{\ell s}}^n\}$ and let $(\mathbf{v}_q^D(\mathbb{F}_{q^{\ell s}}^n, d-1))_{q \in Q}$ be a sequence where each $\mathbf{v}_q^D(\mathbb{F}_{q^{\ell s}}^n, d-1)$ describes the size of the ball of radius $d-1$ in $\mathbb{F}_{q^{\ell s}}^n$. We have

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$$\max \left\{ \liminf_{q \rightarrow +\infty} \left(1 - \mathbf{v}_q^D(\mathbb{F}_{q^{\ell s}}^n, d-1) q^{\ell(k-ns-1)} \right), 0 \right\} \leq \liminf_{q \rightarrow +\infty} \delta_q^D(\mathbb{F}_{q^{\ell s}}^n, q^{\ell k}, \ell, d).$$

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In particular, if $q^{\ell k} \in o\left(\frac{q^{\ell(ns+1)}}{\mathbf{v}_q^D(\mathbb{F}_{q^{\ell s}}^n, d-1)}\right)$ as $q \rightarrow +\infty$, then

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Asymptotic Densities of Extremal Codes

Table: The parameters we consider are q (the field size), ℓ (the index of sublinearity) and s (the \mathbb{F}_{q^ℓ} -dimension of \mathbb{F}_{q^m}).

	$\lim_{q \rightarrow +\infty} \delta_q^D(\mathbb{F}_{q^m}^n, \mathcal{S}, \ell, d)$	$\lim_{\ell \rightarrow +\infty} \delta_q^D(\mathbb{F}_{q^m}^n, \mathcal{S}, \ell, d)$	$\lim_{s \rightarrow +\infty} \delta_q^D(\mathbb{F}_{q^m}^n, \mathcal{S}, \ell, d)$
sublinear MDS	1	1	$\leq (1 + q^{-\ell} \binom{n}{d-1})^{-1}$
non-linear MDS	0 ¹	-	-

¹ A. Gruica, A. Ravagnani; "The Typical Non-Linear Code over Large Alphabets"; 2021 IEEE Information Theory Workshop (ITW)t

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sublinear MDS	1	1	$\leq (1 + q^{-\ell} \binom{n}{d-1})^{-1}$
non-linear MDS	0^1	-	-
sublinear MRD	0^2	1^3	$\leq \left(1 + q^{-\ell} \begin{bmatrix} n \\ d-1 \end{bmatrix}_q\right)^{-1}$
non-linear MRD	0^4	-	-

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- ² A. Gruica, A. Ravagnani; "Common complements of linear subspaces and the sparseness of MRD codes"; SIAM Journal on Applied Algebra and Geometry; 2022
- ³ A.Neri, A.-L. Horlemann-Trautmann, T. Randrianarisoa, J. Rosenthal; "On the genericity of maximum rank distance and Gabidulin codes"; Designs, Codes and Cryptography; 2018
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sublinear MRD	0^2	1^3	$\leq \left(1 + q^{-\ell} \left[\begin{matrix} n \\ d-1 \end{matrix} \right]_q\right)^{-1}$
non-linear MRD	0^4	-	-
sublinear MSRD	various	1^5	$\leq \left(1 + q^{-\ell} \sum_{\substack{(j_1, \dots, j_t) \in \mathbb{N}_0^t \\ j_1 + \dots + j_t = d-1}} \prod_{i=1}^t \left[\begin{matrix} n_i \\ j_i \end{matrix} \right]_q\right)^{-1}$
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⁴ A. Gruica, A. Ravagnani; "The Proportion of (Non-) Linear MRD Codes".

⁵ C.Ott, S. Puchinger, M. Bossert; "Bounds and genericity of sum-rank-metric codes"; 2021 XVII International Symposium "Problems of Redundancy in Information and Control Systems"

- Let $x \in \mathbb{F}_{q^m}^n$ consisting of t blocks $x^{(i)} \in \mathbb{F}_{q^m}^{n_i}$ for $i \in [t]$. The **sum-rank weight** of x is defined as $\omega^{sr,t}(x) := \sum_{i=1}^t \omega^{rk}(x^{(i)})$, where $\omega^{rk}(x^{(i)}) := \dim_{\mathbb{F}_q} \langle x_1^{(i)}, \dots, x_{n_i}^{(i)} \rangle$. For $x, y \in \mathbb{F}_{q^m}^n$ the **sum-rank distance** is defined via $D^{sr,t}(x, y) = \omega^{sr,t}(x - y)$.

¹E. Byrne, H. Gluesing-Luerssen and A. Ravagnani, "Fundamental properties of sum-rank-metric codes," *IEEE Transactions on Information Theory*, **journal** 67, **number** 10, pages 6456–6475, 2021.

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- ▶ $\mathbf{v}_q^{sr,t}(\mathbb{F}_{q^m}^n, r) := \sum_{j=0}^r \sum_{\substack{(j_1, \dots, j_t) \in \mathbb{N}_0^t \\ j_1 + \dots + j_t = j}} \prod_{i=1}^t \begin{bmatrix} n_i \\ j_i \end{bmatrix}_q \prod_{k=0}^{j_i-1} (q^m - q^k)$ denotes the **volume of the sum-rank metric ball** of radius $0 \leq r \leq n$ in $\mathbb{F}_{q^m}^n$.

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Codes in the Sum-Rank Metric¹

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- ▶ **Singleton-type bound** : $|\mathcal{C}| \leq q^{m(n-D(\mathcal{C})+1)}$
- ▶ If $\mathcal{C} \subseteq \mathbb{F}_{q^m}^n$ attains the Singleton-type bound or if $|\mathcal{C}| = 1$, then \mathcal{C} is called **MSRD-** (maximum sum-rank distance) **code**.

¹E. Byrne, H. Gluesing-Luerssen and A. Ravagnani, "Fundamental properties of sum-rank-metric codes," *IEEE Transactions on Information Theory*, **journal** 67, **number** 10, pages 6456–6475, 2021.

Theorem

For $2 \leq d \leq n$ we have

$$\lim_{q \rightarrow +\infty} \delta_q^{\text{sr,t}}(\mathbb{F}_q^n, q^{n-d+1}, 0, d) = 0.$$

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Proof.

Using the estimate

$$\mathbf{v}_q^{\text{sr,t}}(\mathbb{F}_q^n, d-1) \sim q^{(d-1)} \sum_{\substack{(j_1, \dots, j_t) \in \mathbb{N}_0^t, \\ j_1 + \dots + j_t = d-1}} q^{(j_1(n_1 - j_1) + \dots + j_t(n_t - j_t))}$$

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$$\lim_{q \rightarrow +\infty} \frac{q^n}{\mathbf{v}_q^{\text{sr,t}}(\mathbb{F}_q^n, d-1) q^{2(n-d+1)}} \leq \lim_{q \rightarrow +\infty} \frac{1}{q^{n-d+1}} = 0.$$

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Therefore $q^{n-d+1} \in \omega \left(\sqrt{\frac{q^n}{\mathbf{v}_q^{\text{sr,t}}(\mathbb{F}_q^n, d-1)}} \right)$ as $q \rightarrow +\infty$. □

Theorem

Let $m, d, n_1, \dots, n_t \geq 1$ be integers such that $n = \sum_{i=1}^t n_i$ and $1 \leq d \leq n$.

(i) If $d = 1$ or $t = n$ and $n_i = 1$ for all $i \in [n]$ we have

$$\lim_{q \rightarrow +\infty} \delta_q^{sr,t}(\mathbb{F}_{q^m}^n, q^{m(n-d+1)}, 1, d) = 1.$$

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$$\limsup_{q \rightarrow +\infty} \delta_q^{sr,t}(\mathbb{F}_{q^m}^n, q^{m(n-d+1)}, 1, d) \leq \frac{1}{1 + \binom{t-1}{d-2}}.$$

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$$\lim_{q \rightarrow +\infty} \delta_q^{sr,t}(\mathbb{F}_{q^m}^n, q^{m(n-d+1)}, 1, d) = 1.$$

(ii) If $n = t + 1$ we have

$$\limsup_{q \rightarrow +\infty} \delta_q^{sr,t}(\mathbb{F}_{q^m}^n, q^{m(n-d+1)}, 1, d) \leq \frac{1}{1 + \binom{t-1}{d-2}}.$$

(iii) If $d = 2$, $n_{i_1} = \dots = n_{i_r} = 2$ and $n_{i_{r+1}} = \dots = n_{i_t} = 1$ for $r \leq t$ we have

$$\limsup_{q \rightarrow +\infty} \delta_q^{sr,t}(\mathbb{F}_{q^m}^n, q^{m(n-d+1)}, 1, d) \leq \frac{1}{1+r}.$$

Theorem (cont.)

- (iv) In all other cases, that is, if $d \geq 3$ and $n \geq t + 2$ or $d = 2$ and $\max\{n_1, \dots, n_t\} \geq 3$, we have

$$\lim_{q \rightarrow +\infty} \delta_q^{sr,t}(\mathbb{F}_{q^m}^n, q^{m(n-d+1)}, 1, d) = 0.$$

Theorem (cont.)

- (iv) In all other cases, that is, if $d \geq 3$ and $n \geq t + 2$ or $d = 2$ and $\max\{n_1, \dots, n_t\} \geq 3$, we have

$$\lim_{q \rightarrow +\infty} \delta_q^{sr,t}(\mathbb{F}_{q^m}^n, q^{m(n-d+1)}, 1, d) = 0.$$

Proof.

Using the estimate

$$\mathbf{v}_q^{sr,t}(\mathbb{F}_{q^m}^n, d-1) \sim q^{m(d-1)} \sum_{\substack{(j_1, \dots, j_t) \in \mathbb{N}_0^t, \\ j_1 + \dots + j_t = d-1}} q^{(j_1(n_1-j_1) + \dots + j_t(n_t-j_t))}$$

gives us

Theorem (cont.)

- (iv) In all other cases, that is, if $d \geq 3$ and $n \geq t + 2$ or $d = 2$ and $\max\{n_1, \dots, n_t\} \geq 3$, we have

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gives us

$$\lim_{q \rightarrow \infty} \frac{q^{m(n-d+1)} \cdot \mathbf{v}_q^{sr,t}(\mathbb{F}_{q^m}^n, d-1)}{q^{(nm+1)}} = \lim_{q \rightarrow \infty} \frac{1}{q} \left(\sum_{\substack{(j_1, \dots, j_t) \in \mathbb{N}_0^t, \\ j_1 + \dots + j_t = d-1}} q^{(j_1(n_1-j_1) + \dots + j_t(n_t-j_t))} \right).$$

Theorem (cont.)

- (iv) In all other cases, that is, if $d \geq 3$ and $n \geq t + 2$ or $d = 2$ and $\max\{n_1, \dots, n_t\} \geq 3$, we have

$$\lim_{q \rightarrow +\infty} \delta_q^{sr,t}(\mathbb{F}_{q^m}^n, q^{m(n-d+1)}, 1, d) = 0.$$

Proof.

Using the estimate

$$\mathbf{v}_q^{sr,t}(\mathbb{F}_{q^m}^n, d-1) \sim q^{m(d-1)} \sum_{\substack{(j_1, \dots, j_t) \in \mathbb{N}_0^t, \\ j_1 + \dots + j_t = d-1}} q^{(j_1(n_1 - j_1) + \dots + j_t(n_t - j_t))}$$

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Now analyze the four cases separately by using combinatorial arguments.

Thank you for your attention!